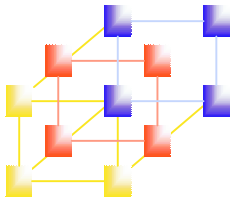
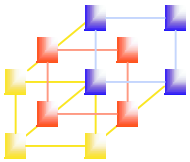


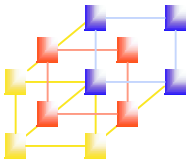
Chapter 4 Macro Processors

-- Basic Macro Processor Functions



Introduction

- A macro instruction (macro) is a notational convenience for the programmer
  - It allows the programmer to write shorthand version of a program (module programming)
- The macro processor replaces each macro instruction with the corresponding group of source language statements (*expanding*)
  - Normally, it performs no analysis of the text it handles.
  - It does not concern the meaning of the involved statements during macro expansion.
- The design of a macro processor generally is *machine independent!*

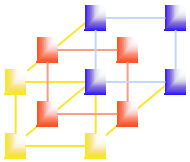


Basic macro processor functions

- Two new assembler directives are used in macro definition
  - **MACRO:** identify the beginning of a macro definition
  - **MEND:** identify the end of a macro definition
- Prototype for the macro
  - Each parameter begins with '&'

```
name    MACRO    parameters
                    :
                 body
                    :
            MEND
```

  - Body: the statements that will be generated as the expansion of the macro.



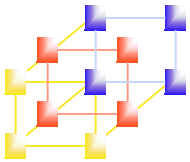
Macro expansion

Source

```
M1      MACRO   &D1, &D2
        STA     &D1
        STB     &D2
        MEND
    .
M1 DATA1, DATA2
    .
M1 DATA4, DATA3
```

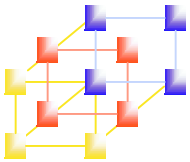
Expanded source

```
    .
    .
    .
        STA     DATA1
        STB     DATA2
    .
        STA     DATA4
        STB     DATA3
    .
```



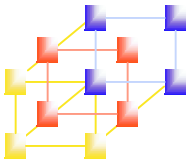
Example of macro definition Figure 4.1, pp. 178

```
5    COPY      START   0                 COPY FILE FROM INPUT TO OUTPUT
10   RDBUFF    MACRO   &INDEV, &BUFADR &RECLTH
15   .
20   .         MACRO TO READ RECORD INTO BUFFER
25   .
30             CLEAR   X                 CLEAR LOOP COUNTER
35             CLEAR   A
40             CLEAR   S
45             +LDT    #4096             SET MAXINUM RECORD LENTH
50             TD      =X'&INDEV'        TEST INPUT DEVICE
55             JEQ     *-3               LOOP UNTIL READY
60             RD      =X'&INDEV'        READ CHARACTER INTO REG A
65             COMPR   A , S             TEST FOR END OF RECORD
70             JEQ     *+11              EXIT LOOP IF EOR
75             STCH    &BUFADR, X        STORE CHARACTER IN BUFFER
80             TIXR    T                 LOOP UNLESS MAXIMUN LENGTH
85             JLT     *-19              HAS BEEN RECARD
90             STX     &RECLTH           SAVE RECORD LENGTH
95             MEND
```



Macro invocation

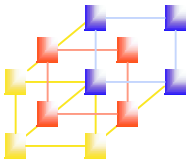
- A macro invocation statement (a macro call) gives the name of the macro instruction being invoked and the arguments to be used in expanding the macro.
  - `macro_name  p1, p2, …`
- Difference between macro call and procedure call
  - Macro call: statements of the macro body are expanded each time the macro is invoked.
  - Procedure call: statements of the subroutine appear only one, regardless of how many times the subroutine is called.
- Question
  - How does a programmer decide to use macro calls or procedure calls?
    - `From the viewpoint of a programmer`
    - `From the viewpoint of the CPU`



Exchange the values of two variables

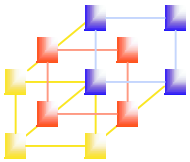
```
void exchange(int a, int b) {
  int temp;
  temp = a;
  a = b;
  b = temp;
  }

main() {
  int i=1, j=3;
  printf("BEFORE - %d %d\n", i, j);
  exchange(i, j);
  printf("AFTER - %d %d\n", i, j);
}
```

What's the result?



Pass by Reference

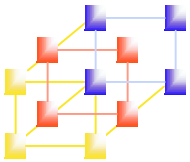
```
void exchange(int *p1, int *p2) {
  int temp;
  temp = *p1;
  *p1 = *p2;
  *p2 = temp;
  }

main() {
  int i=1, j=3;
  printf("BEFORE - %d %d\n", i, j);
  exchange(&i, &j);
  printf("AFTER - %d %d\n", i, j);
}
```



12 Lines of Assembly Code

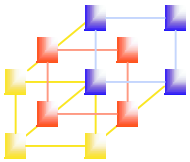
```
. Subroutine EXCH
EXCH      LDA      @P1
          STA       TEMP
          LDA      @P2
          STA      @P1
          LDA       TEMP
          STA      @P2
          RSUB
P1        RESW      1
P2        RESW      1
TEMP      RESW      1
```

```
MAIN      LDA      #1
          STA       I
          LDA      #3
          STA       J
. Call a subroutine
          LDA      #I
          STA       P1
          LDA      #J
          STA       P2
          JSUB      EXCH
I         RESW      1
J         RESW      1
          END       MAIN
```



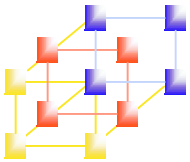
Swap two variables by macro

```
#define swap(i,j) { int temp; temp=i; i=j; j=temp; }

main() {
   int i=1, j=3;
   printf("BEFORE - %d %d\n", i, j);
   swap(i,j);
   printf("AFTER - %d %d\n", i, j);
}
```



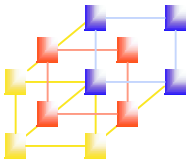
6 Lines of Assembly Code

```
MAIN      LDA      #1
          STA       I
          LDA      #3
          STA       J
. Invoke a macro
          LDA       I
          STA       TEMP
          LDA       J
          STA       I
          LDA       TEMP
          STA       J
I         RESW      1
J         RESW      1
TEMP      RESW      1
          END       MAIN
```



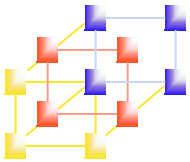
Macro expansion

- Each macro invocation statement will be expanded into the statements that form the body of the macro.
- Arguments from the macro invocation are substituted for the parameters in the macro prototype (according to their positions).
  - In the definition of macro: parameter
  - In the macro invocation: argument
- Comment lines within the macro body will be deleted.
- Macro invocation statement itself has been included as a comment line.
- The label on the macro invocation statement has been retained as a label on the first statement generated in the macro expansion.
  - We can use a macro instruction in exactly the same way as an assembler language mnemonic.



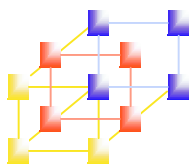
Example of macro invocation Figure 4.1, pp. 178

```
170  .                    MAIN PROGRAM
175  .
180       FIRST    STL     RETADR              SAVE RETURN ADDRESS
190       CLOOP    RDBUFF  F1,BUFFER,LENGTH    READ RECORD INTO BUFFER
195                LDA     LENGTH              TEST FOR END OF FILE
200                COMP    #0
205                JEQ     ENDFIL              EXIT IF EOF FOUND
210                WRBUFF  05,BUFFER,LENGTH    WRITE OUTPUT RECORD
215                J       CLOOP               LOOP
220       ENDFIL   WRBUFF  05,EOF,THREE        INSERT EOF MARKER
225                J       @RETADR
230       EOF      BYTE    C'EOF'
235       THREE    WORD    3
240       RETADR   RESW    1
245       LENGTH   RESW    1                   LENGTH OF RECORD
250       BUFFER   RESB    4096                4096-BYTE BUFFER AREA
255                END     FIRST
```



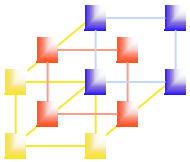
Example of macro expansion
Figure 4.2, pp. 179

```
5       COPY    START   0               COPY FILE FROM INPUT TO OUTPUT
180     FIRST   STL     RETADR          SAVE RETURN ADDRESS
190     .CLOOP  RDBUFF  F1,BUFFER,LENGTH  READ RECORD INTO BUFFER
190a    CLOOP   CLEAR   X               CLEAR LOOP COUNTER
190b            CLEAR   A
190c            CLEAR   S
190d            +LDT    #4096           SET MAXIMUN RECORD LENGTH
190e            TD      =X'F1'          TEST INPUT DEVICE
190f            JEQ     *-3             LOOP UNTIL READY
190g            RD      =X'F1'          TEST FOR END OF RECORD
190h            COMPR   A, S            TEST FOR END OF RECORD
190i            JEQ     *+11            EXIT LOOP IF EOR
190j            STCH    BUFFER, X       STORE CHARACTER IN BUFFER
190k            TIXR    T               LOOP UNLESS MAXIMUN LENGTH
190l            JLT     *-19            HAS BEEN REACHED
190M            STX     LENGTH          SAVE RECORD LENGTH
```



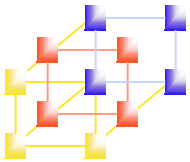
Example of macro expansion
Figure 4.2, pp. 179

```
195                 LDA     LENGTH          TEST FOR END OF FILE
200                 COMP    #0
205                 JEQ     ENDFIL          EXIT IF EOF FOUND
210                 WRBUFF  05,BUFFER,LENGTH WRITE OUTPUT RECORD
210a                CLEAR   X               CLEAR LOOP COUNTER
210b                LDT     LENGTH
210c                LDCH    BUFFER,X        GET CHARACTER FROM BUFFER
210d                TD      =X'05'          TEST OUTPUT DEVICE
210e                JEQ     *-3             LOOP UNTIL READY
210f                WD      =X'05'          WRITE CHARACTER
210g                TIXR    T               LOOP UNTIL ALL CHARACTERS
210h                JLT     *-14            HAVE BEEN WRITTEN
215                 J       CLOOP           LOOP
220         .ENDFIL WRBUFF  05,EOF,THREE    INSERT EOF MARKER
```



Example of macro expansion Figure 4.2, pp. 179

```
220a    ENDFIL  CLEAR   X       CLEAR LOOP COUNTER
220b            LDT     THREE
220c            LDCH    EOF,X   GET CHARACTER FROM BUFFER
220d            TD      =X'05'  TEST OUTPUT DEVICE
220e            JEQ     *-3     LOOP UNTIL READY
220f            WD      =X'05'  WRITE CHARACTER
220g            TIXR    T       LOOP UNTIL ALL CHARACTERS
220h            JLT     *-14    HAVE BEEN WRITTEN
225             J       @RETADR
230     EOF     BYTE    C'EOF'
235     THREE   WORD    3
240     RETADR  RESW    1
245     LENGTH  RESW    1
250     BUFFER  RESB    4096    4096-BYTE BUFFER AREA
255             END     FIRST
```

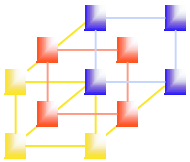


No label in the macro body

- Problem of the label in the body of macro:
  - If the same macro is expanded multiple times at different places in the program …
  - There will be *duplicate labels*, which will be treated as errors by the assembler.
- Solutions:
  - Do not use labels in the body of macro.
  - Explicitly use PC-relative addressing instead.
    - Ex, in RDBUFF and WRBUFF macros,

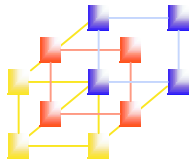
```
JEQ    *+11
JLT    *-14
```

    - It is inconvenient and error-prone.
  - The way of avoiding such error-prone method will be discussed in Section 4.2.2



Two-pass macro processor

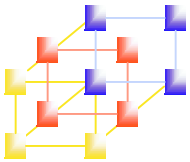
- You may design a two-pass macro processor
  - Pass 1:
    - `Process all macro definitions`
  - Pass 2:
    - `Expand all macro invocation statements`
- However, one-pass may be enough
  - Because all macros would have to be defined during the first pass before any macro invocations were expanded.
    - `The definition of a macro must appear before any statements that invoke that macro.`
  - Moreover, the body of one macro can contain definitions of other macros.



Example of recursive macro definition Figure 4.3, pp.182

- MACROS (for SIC)
  - Contains the definitions of RDBUFF and WRBUFF written in SIC instructions.

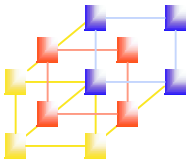
```
1        MACROS   MACOR      {Defines SIC standard version macros}
2        RDBUFF   MACRO      &INDEV,&BUFADR,&RECLTH
                  .
                  .          {SIC standard version}
                  .
3                  MEND      {End of RDBUFF}
4        WRBUFF   MACRO      &OUTDEV,&BUFADR,&RECLTH
                  .
                  .          {SIC standard version}
5                  MEND      {End of WRBUFF}
                  .
                  .
                  .
6                  MEND      {End of MACROS}
```



Example of recursive macro definition Figure 4.3, pp.182

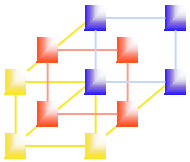
- MACROX (for SIC/XE)
  - Contains the definitions of RDBUFF and WRBUFF written in SIC/XE instructions.

```
1          MACROX   MACRO      {Defines SIC/XE macros}
2          RDBUFF   MACRO      &INDEV,&BUFADR,&RECLTH
                    .
                    .          {SIC/XE version}
                    .
3                   MEND       {End of RDBUFF}
4          WRBUFF   MACRO      &OUTDEV,&BUFADR,&RECLTH
                    .
                    .          {SIC/XE version}
                    .
5                   MEND       {End of WRBUFF}
                    .
                    .
6                   MEND       {End of MACROX}
```



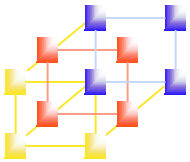
Example of macro definitions

- A program that is to be run on SIC system could invoke MACROS whereas a program to be run on SIC/XE can invoke MACROX.
- However, defining MACROS or MACROX does not define RDBUFF and WRBUFF. These definitions are processed only when an invocation of MACROS or MACROX is expanded.



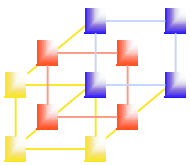
One-pass macro processor

- A one-pass macro processor that alternate between *macro definition* and *macro expansion* in a recursive way is able to handle recursive macro definition.
- Restriction
  - The definition of a macro must appear in the source program before any statements that invoke that macro.
  - This restriction does not create any real inconvenience.

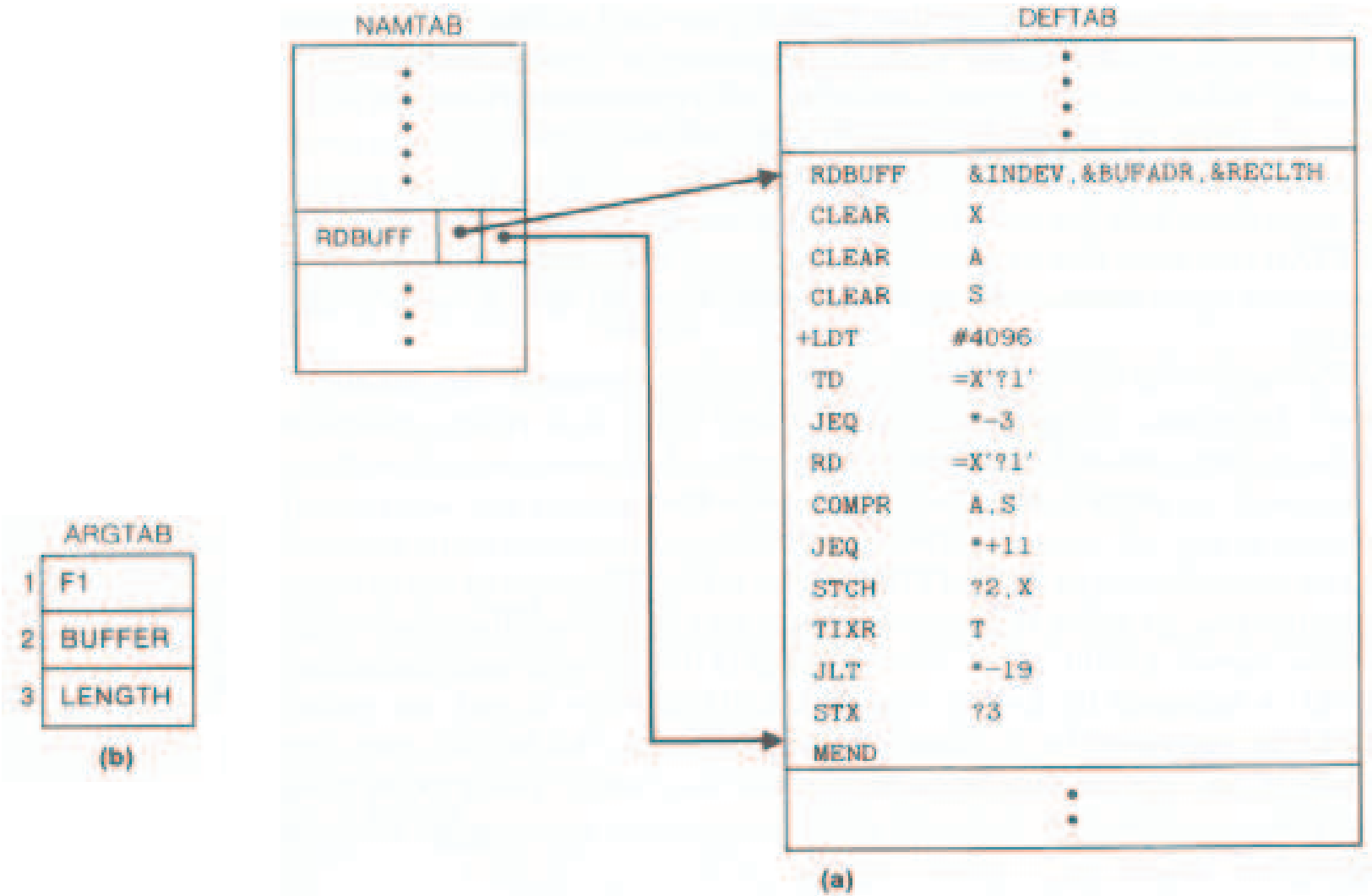


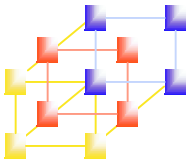
Data structures for one-pass macro processor

- DEFTAB (definition table)
  - Stores the macro definition including *macro prototype* and *macro body*
  - Comment lines are omitted.
  - References to the macro instruction parameters are converted to a positional notation for efficiency in substituting arguments.
- NAMTAB
  - Stores macro names
  - Serves as an index to DEFTAB
    - `Pointers to the beginning and the end of the macro definition (DEFTAB)`
- ARGTAB
  - Stores the arguments of macro invocation according to their positions in the argument list
  - As the macro is expanded, arguments from ARGTAB are substituted for the corresponding parameters in the macro body.



Data structures





Algorithm

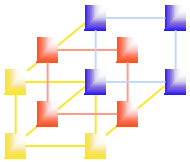
Procedure GETLINE
If EXPANDING **then**
get the next line to be processed from DEFTAB
Else
read next line from input file

MAIN program
- Iterations of
  - GETLINE
  - PROCESSLINE

Procedure PROCESSLINE
- DEFINE
- EXPAND
- Output source line

Procedure EXPAND
Set up the argument values in ARGTAB
Expand a macro invocation statement (like in MAIN procedure)
- Iterations of
  - GETLINE
  - PROCESSLINE

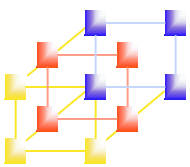
Procedure DEFINE
Make appropriate entries in DEFTAB and NAMTAB



Algorithm
Figure 4.5, pp. 184

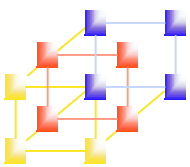
```
begin {macro processor}
        EXPANDINF := FALSE
        while OPCODE ≠ 'END' do
                begin
                        GETLINE
                        PROCESSLINE
                end {while}
end {macro processor}
```

```
Procedure PROCESSLINE
        begin
          search MAMTAB for OPCODE
          if found then
                EXPAND
          else if OPCODE = 'MACRO' then
                DEFINE
          else write source line to expanded file
        end {PRCOESSOR}
```



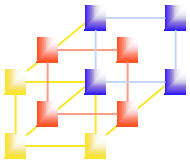
Algorithm
Figure 4.5, pp. 185

```
Procedure DEFINE
        begin
            enter macro name into NAMTAB
            enter macro prototype into DEFTAB
            LEVEL   :- 1
            while LEVEL > do
                begin
                    GETLINE
                    if this is not a comment line then
                      begin
                        substitute positional notation for parameters
                        enter line into DEFTAB
                        if OPCODE = 'MACRO' then
                            LEVEL := LEVEL +1
                        else if OPCODE = 'MEND' then
                            LEVEL := LEVEL – 1
                      end {if not comment}
                end {while}
             store in NAMTAB pointers to beginning and end of definition
        end {DEFINE}
```



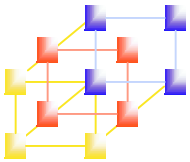
Algorithm
Figure 4.5, pp. 185

```
Procedure EXPAND
    begin
            EXPANDING := TRUE
            get first line of macro definition {prototype} from DEFTAB
            set up arguments from macro invocation in ARGTAB
            while macro invocation to expanded file as a comment
            while not end of macro definition do
                begin
                        GETLINE
                        PROCESSLINE
                end {while}
            EXPANDING := FALSE
    end {EXPAND}

Procedure GETLINE
    begin
            if EXPANDING then
                begin
                    get next line of macro definition from DEFTAB
                    substitute arguments from ARGTAB for positional notation
                end {if}
            else
                read next line from input file
    end {GETLINE}
```



Handling nested macro definition

- In DEFINE procedure
  - When a macro definition is being entered into DEFTAB, the normal approach is to continue until an MEND directive is reached.
  - This would not work for nested macro definition because the first MEND encountered in the inner macro will terminate the whole macro definition process.
  - To solve this problem, a counter LEVEL is used to keep track of the level of macro definitions.
    - `Increase LEVEL by 1 each time a MACRO directive is read.`
    - `Decrease LEVEL by 1 each time a MEND directive is read.`
    - `A MEND terminates the whole macro definition process when LEVEL reaches 0.`
    - `This process is very much like matching left and right parentheses when scanning an arithmetic expression.`



Comparison of Macro Processors Design

- One-pass algorithm
  - Every macro must be defined before it is called
  - One-pass processor can alternate between macro definition and macro expansion
  - Nested macro definitions are allowed but nested calls are not
- Two-pass algorithm
  - Pass1: Recognize macro definitions
  - Pass2: Recognize macro calls
  - Nested macro definitions are not allowed